

Linux generic IRQ handling

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by Thomas Gleixner and Ingo Molnar

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Chapter 1. Introduction

The generic interrupt handling layer is designed to provide a complete abstraction of interrupt handling for device drivers. It is able to handle all the different types of interrupt controller hardware. Device drivers use generic API functions to request, enable, disable and free interrupts. The drivers do not have to know anything about interrupt hardware details, so they can be used on different platforms without code changes.

This documentation is provided to developers who want to implement an interrupt subsystem based for their architecture, with the help of the generic IRQ handling layer.

Chapter 2. Rationale

The original implementation of interrupt handling in Linux is using the `__do_IRQ()` super-handler, which is able to deal with every type of interrupt logic.

Originally, Russell King identified different types of handlers to build a quite universal set for the ARM interrupt handler implementation in Linux 2.5/2.6. He distinguished between:

- Level type
- Edge type
- Simple type

In the SMP world of the `__do_IRQ()` super-handler another type was identified:

- Per CPU type

This split implementation of highlevel IRQ handlers allows us to optimize the flow of the interrupt handling for each specific interrupt type. This reduces complexity in that particular codepath and allows the optimized handling of a given type.

The original general IRQ implementation used `hw_interrupt_type` structures and their `->ack()`, `->end()` [etc.] callbacks to differentiate the flow control in the super-handler. This leads to a mix of flow logic and lowlevel hardware logic, and it also leads to unnecessary code duplication: for example in i386, there is a `ioapic_level_irq` and a `ioapic_edge_irq` irq-type which share many of the lowlevel details but have different flow handling.

A more natural abstraction is the clean separation of the 'irq flow' and the 'chip details'.

Analysing a couple of architecture's IRQ subsystem implementations reveals that most of them can use a generic set of 'irq flow' methods and only need to add the chip level specific code. The separation is also valuable for (sub)architectures which need specific quirks in the irq flow itself but not in the chip-details - and thus provides a more transparent IRQ subsystem design.

Each interrupt descriptor is assigned its own highlevel flow handler, which is normally one of the generic implementations. (This highlevel flow handler implementation also makes it simple to provide demultiplexing handlers which can be found in embedded platforms on various architectures.)

The separation makes the generic interrupt handling layer more flexible and extensible. For example, an (sub)architecture can use a generic irq-flow

Chapter 2. Rationale

implementation for 'level type' interrupts and add a (sub)architecture specific 'edge type' implementation.

To make the transition to the new model easier and prevent the breakage of existing implementations, the `__do_IRQ()` super-handler is still available. This leads to a kind of duality for the time being. Over time the new model should be used in more and more architectures, as it enables smaller and cleaner IRQ subsystems.

Chapter 3. Known Bugs And Assumptions

None (knock on wood).

Chapter 3. Known Bugs And Assumptions

Chapter 4. Abstraction layers

There are three main levels of abstraction in the interrupt code:

1. Highlevel driver API
2. Highlevel IRQ flow handlers
3. Chiplevel hardware encapsulation

4.1. Interrupt control flow

Each interrupt is described by an interrupt descriptor structure `irq_desc`. The interrupt is referenced by an 'unsigned int' numeric value which selects the corresponding interrupt description structure in the descriptor structures array. The descriptor structure contains status information and pointers to the interrupt flow method and the interrupt chip structure which are assigned to this interrupt.

Whenever an interrupt triggers, the lowlevel arch code calls into the generic interrupt code by calling `desc->handle_irq()`. This highlevel IRQ handling function only uses `desc->chip` primitives referenced by the assigned chip descriptor structure.

4.2. Highlevel Driver API

The highlevel Driver API consists of following functions:

- `request_irq()`
- `free_irq()`
- `disable_irq()`
- `enable_irq()`
- `disable_irq_nosync()` (SMP only)
- `synchronize_irq()` (SMP only)
- `set_irq_type()`
- `set_irq_wake()`
- `set_irq_data()`

- `set_irq_chip()`
- `set_irq_chip_data()`

See the autogenerated function documentation for details.

4.3. Highlevel IRQ flow handlers

The generic layer provides a set of pre-defined irq-flow methods:

- `handle_level_irq`
- `handle_edge_irq`
- `handle_simple_irq`
- `handle_percpu_irq`

The interrupt flow handlers (either predefined or architecture specific) are assigned to specific interrupts by the architecture either during bootup or during device initialization.

4.3.1. Default flow implementations

4.3.1.1. Helper functions

The helper functions call the chip primitives and are used by the default flow implementations. The following helper functions are implemented (simplified excerpt):

```
default_enable(irq)
{
    desc->chip->unmask(irq);
}

default_disable(irq)
{
    if (!delay_disable(irq))
        desc->chip->mask(irq);
}

default_ack(irq)
{
    chip->ack(irq);
}
```

```
default_mask_ack(irq)
{
    if (chip->mask_ack) {
        chip->mask_ack(irq);
    } else {
        chip->mask(irq);
        chip->ack(irq);
    }
}

noop(irq)
{
}
```

4.3.2. Default flow handler implementations

4.3.2.1. Default Level IRQ flow handler

`handle_level_irq` provides a generic implementation for level-triggered interrupts.

The following control flow is implemented (simplified excerpt):

```
desc->chip->start();
handle_IRQ_event(desc->action);
desc->chip->end();
```

4.3.2.2. Default Edge IRQ flow handler

`handle_edge_irq` provides a generic implementation for edge-triggered interrupts.

The following control flow is implemented (simplified excerpt):

```
if (desc->status & running) {
    desc->chip->hold();
    desc->status |= pending | masked;
    return;
}
```

```
desc->chip->start();
desc->status |= running;
do {
    if (desc->status & masked)
        desc->chip->enable();
    desc->status &= ~pending;
    handle_IRQ_event(desc->action);
} while (status & pending);
desc->status &= ~running;
desc->chip->end();
```

4.3.2.3. Default simple IRQ flow handler

`handle_simple_irq` provides a generic implementation for simple interrupts.

Note: The simple flow handler does not call any handler/chip primitives.

The following control flow is implemented (simplified excerpt):

```
handle_IRQ_event(desc->action);
```

4.3.2.4. Default per CPU flow handler

`handle_percpu_irq` provides a generic implementation for per CPU interrupts.

Per CPU interrupts are only available on SMP and the handler provides a simplified version without locking.

The following control flow is implemented (simplified excerpt):

```
desc->chip->start();
handle_IRQ_event(desc->action);
desc->chip->end();
```

4.3.3. Quirks and optimizations

The generic functions are intended for 'clean' architectures and chips, which have no platform-specific IRQ handling quirks. If an architecture needs to implement quirks on the 'flow' level then it can do so by overriding the highlevel irq-flow handler.

4.3.4. Delayed interrupt disable

This per interrupt selectable feature, which was introduced by Russell King in the ARM interrupt implementation, does not mask an interrupt at the hardware level when `disable_irq()` is called. The interrupt is kept enabled and is masked in the flow handler when an interrupt event happens. This prevents losing edge interrupts on hardware which does not store an edge interrupt event while the interrupt is disabled at the hardware level. When an interrupt arrives while the `IRQ_DISABLED` flag is set, then the interrupt is masked at the hardware level and the `IRQ_PENDING` bit is set. When the interrupt is re-enabled by `enable_irq()` the pending bit is checked and if it is set, the interrupt is resent either via hardware or by a software resend mechanism. (It's necessary to enable `CONFIG_HARDIRQS_SW_RESEND` when you want to use the delayed interrupt disable feature and your hardware is not capable of retriggering an interrupt.) The delayed interrupt disable can be runtime enabled, per interrupt, by setting the `IRQ_DELAYED_DISABLE` flag in the `irq_desc` status field.

4.4. Chiplevel hardware encapsulation

The chip level hardware descriptor structure `irq_chip` contains all the direct chip relevant functions, which can be utilized by the irq flow implementations.

- `ack()`
- `mask_ack()` - Optional, recommended for performance
- `mask()`
- `unmask()`
- `retrigger()` - Optional
- `set_type()` - Optional
- `set_wake()` - Optional

Chapter 4. Abstraction layers

These primitives are strictly intended to mean what they say: ack means ACK, masking means masking of an IRQ line, etc. It is up to the flow handler(s) to use these basic units of lowlevel functionality.

Chapter 5. `__do_IRQ` entry point

The original implementation `__do_IRQ()` is an alternative entry point for all types of interrupts.

This handler turned out to be not suitable for all interrupt hardware and was therefore reimplemented with split functionality for `edge/level/simple/percpu` interrupts. This is not only a functional optimization. It also shortens code paths for interrupts.

To make use of the split implementation, replace the call to `__do_IRQ` by a call to `desc->handle_irq()` and associate the appropriate handler function to `desc->handle_irq()`. In most cases the generic handler implementations should be sufficient.

Chapter 6. Locking on SMP

The locking of chip registers is up to the architecture that defines the chip primitives. There is a `chip->lock` field that can be used for serialization, but the generic layer does not touch it. The `per-irq` structure is protected via `desc->lock`, by the generic layer.

Chapter 7. Structures

This chapter contains the autogenerated documentation of the structures which are used in the generic IRQ layer.

struct irq_chip

LINUX

Kernel Hackers Manual July 2010

Name

struct irq_chip — hardware interrupt chip descriptor

Synopsis

```
struct irq_chip {
    const char * name;
    unsigned int (* startup) (unsigned int irq);
    void (* shutdown) (unsigned int irq);
    void (* enable) (unsigned int irq);
    void (* disable) (unsigned int irq);
    void (* ack) (unsigned int irq);
    void (* mask) (unsigned int irq);
    void (* mask_ack) (unsigned int irq);
    void (* unmask) (unsigned int irq);
    void (* eoi) (unsigned int irq);
    void (* end) (unsigned int irq);
    int (* set_affinity) (unsigned int irq, const struct cpumask *dest);
    int (* retrigger) (unsigned int irq);
    int (* set_type) (unsigned int irq, unsigned int flow_type);
    int (* set_wake) (unsigned int irq, unsigned int on);
    void (* bus_lock) (unsigned int irq);
    void (* bus_sync_unlock) (unsigned int irq);
#ifdef CONFIG_IRQ_RELEASE_METHOD
    void (* release) (unsigned int irq, void *dev_id);
#endif
    const char * typename;
};
```

Members

name

name for /proc/interrupts

startup

start up the interrupt (defaults to ->enable if NULL)

shutdown

shut down the interrupt (defaults to ->disable if NULL)

enable

enable the interrupt (defaults to chip->unmask if NULL)

disable

disable the interrupt

ack

start of a new interrupt

mask

mask an interrupt source

mask_ack

ack and mask an interrupt source

unmask

unmask an interrupt source

eoi

end of interrupt - chip level

end

end of interrupt - flow level

set_affinity

set the CPU affinity on SMP machines

retrigger

resend an IRQ to the CPU

<code>set_type</code>	set the flow type (<code>IRQ_TYPE_LEVEL/etc.</code>) of an IRQ
<code>set_wake</code>	enable/disable power-management wake-on of an IRQ
<code>bus_lock</code>	function to lock access to slow bus (i2c) chips
<code>bus_sync_unlock</code>	function to sync and unlock slow bus (i2c) chips
<code>release</code>	release function solely used by UML
<code>typename</code>	obsoleted by name, kept as migration helper

struct irq_desc

LINUX

Kernel Hackers Manual July 2010

Name

`struct irq_desc` — interrupt descriptor

Synopsis

```
struct irq_desc {
    unsigned int irq;
    struct timer_rand_state * timer_rand_state;
    unsigned int * kstat_irqs;
#ifdef CONFIG_INTR_REMAP
    struct irq_2_iommu * irq_2_iommu;
#endif
    irq_flow_handler_t handle_irq;
}
```

Chapter 7. Structures

```
struct irq_chip * chip;
struct msi_desc * msi_desc;
void * handler_data;
void * chip_data;
struct irqaction * action;
unsigned int status;
unsigned int depth;
unsigned int wake_depth;
unsigned int irq_count;
unsigned long last_unhandled;
unsigned int irqs_unhandled;
raw_spinlock_t lock;
#ifdef CONFIG_SMP
    cpumask_var_t affinity;
    unsigned int node;
#endif
#ifdef CONFIG_GENERIC_PENDING_IRQ
    cpumask_var_t pending_mask;
#endif
#endif
    atomic_t threads_active;
    wait_queue_head_t wait_for_threads;
#ifdef CONFIG_PROC_FS
    struct proc_dir_entry * dir;
#endif
    const char * name;
};
```

Members

irq

interrupt number for this descriptor

timer_rand_state

pointer to timer rand state struct

kstat_irqs

irq stats per cpu

irq_2_iommu

iommu with this irq

`handle_irq`
highlevel irq-events handler [if NULL, `__do_IRQ`]

`chip`
low level interrupt hardware access

`msi_desc`
MSI descriptor

`handler_data`
per-IRQ data for the `irq_chip` methods

`chip_data`
platform-specific per-chip private data for the chip methods, to allow shared chip implementations

`action`
the irq action chain

`status`
status information

`depth`
disable-depth, for nested `irq_disable` calls

`wake_depth`
enable depth, for multiple `set_irq_wake` callers

`irq_count`
stats field to detect stalled irqs

`last_unhandled`
aging timer for unhandled count

`irqs_unhandled`
stats field for spurious unhandled interrupts

`lock`
locking for SMP

affinity

IRQ affinity on SMP

node

node index useful for balancing

pending_mask

pending rebalanced interrupts

threads_active

number of irqaction threads currently running

wait_for_threads

wait queue for sync_irq to wait for threaded handlers

dir

/proc/irq/ procfs entry

name

flow handler name for /proc/interrupts output

alloc_desc_masks

LINUX

Kernel Hackers Manual July 2010

Name

`alloc_desc_masks` — allocate cpumasks for `irq_desc`

Synopsis

```
bool alloc_desc_masks (struct irq_desc * desc, int node, bool boot);
```

Arguments

desc

pointer to `irq_desc` struct

node

node which will be handling the cpumasks

boot

true if need bootmem

Description

Allocates affinity and `pending_mask` cpumask if required. Returns true if successful (or not required).

init_copy_desc_masks

LINUX

Kernel Hackers Manual July 2010

Name

`init_copy_desc_masks` — copy cpumasks for `irq_desc`

Synopsis

```
void init_copy_desc_masks (struct irq_desc * old_desc, struct  
irq_desc * new_desc);
```

Arguments

old_desc

pointer to old irq_desc struct

new_desc

pointer to new irq_desc struct

Description

Insures affinity and pending_masks are copied to new irq_desc. If !CONFIG_CPUMASKS_OFFSTACK the cpumasks are embedded in the irq_desc struct so the copy is redundant.

struct irqaction

LINUX

Kernel Hackers Manual July 2010

Name

struct irqaction — per interrupt action descriptor

Synopsis

```
struct irqaction {
    irq_handler_t handler;
    unsigned long flags;
    const char * name;
    void * dev_id;
    struct irqaction * next;
    int irq;
    struct proc_dir_entry * dir;
    irq_handler_t thread_fn;
    struct task_struct * thread;
    unsigned long thread_flags;
```

```
};
```

Members

handler

interrupt handler function

flags

flags (see `IRQF_*` above)

name

name of the device

dev_id

cookie to identify the device

next

pointer to the next irqaction for shared interrupts

irq

interrupt number

dir

pointer to the `proc/irq/NN/name` entry

thread_fn

interrupt handler function for threaded interrupts

thread

thread pointer for threaded interrupts

thread_flags

flags related to *thread*

Chapter 8. Public Functions Provided

This chapter contains the autogenerated documentation of the kernel API functions which are exported.

synchronize_irq

LINUX

Kernel Hackers Manual July 2010

Name

`synchronize_irq` — wait for pending IRQ handlers (on other CPUs)

Synopsis

```
void synchronize_irq (unsigned int irq);
```

Arguments

irq

interrupt number to wait for

Description

This function waits for any pending IRQ handlers for this interrupt to complete before returning. If you use this function while holding a resource the IRQ handler may need you will deadlock.

This function may be called - with care - from IRQ context.

disable_irq_nosync

LINUX

Kernel Hackers Manual July 2010

Name

`disable_irq_nosync` — disable an irq without waiting

Synopsis

```
void disable_irq_nosync (unsigned int irq);
```

Arguments

irq

Interrupt to disable

Description

Disable the selected interrupt line. Disables and Enables are nested. Unlike `disable_irq`, this function does not ensure existing instances of the IRQ handler have completed before returning.

This function may be called from IRQ context.

disable_irq

LINUX

Name

`disable_irq` — disable an irq and wait for completion

Synopsis

```
void disable_irq (unsigned int irq);
```

Arguments

irq

Interrupt to disable

Description

Disable the selected interrupt line. Enables and Disables are nested. This function waits for any pending IRQ handlers for this interrupt to complete before returning. If you use this function while holding a resource the IRQ handler may need you will deadlock.

This function may be called - with care - from IRQ context.

`enable_irq`

LINUX

Name

`enable_irq` — enable handling of an irq

Synopsis

```
void enable_irq (unsigned int irq);
```

Arguments

irq

Interrupt to enable

Description

Undoes the effect of one call to `disable_irq`. If this matches the last disable, processing of interrupts on this IRQ line is re-enabled.

This function may be called from IRQ context only when `desc->chip->bus_lock` and `desc->chip->bus_sync_unlock` are NULL !

set_irq_wake

LINUX

Kernel Hackers Manual July 2010

Name

`set_irq_wake` — control irq power management wakeup

Synopsis

```
int set_irq_wake (unsigned int irq, unsigned int on);
```

Arguments

irq

interrupt to control

on

enable/disable power management wakeup

Description

Enable/disable power management wakeup mode, which is disabled by default. Enables and disables must match, just as they match for non-wakeup mode support.

Wakeup mode lets this IRQ wake the system from sleep states like “suspend to RAM”.

setup_irq

LINUX

Kernel Hackers Manual July 2010

Name

`setup_irq` — setup an interrupt

Synopsis

```
int setup_irq (unsigned int irq, struct irqaction * act);
```

Arguments

irq

Interrupt line to setup

act

irqaction for the interrupt

Description

Used to statically setup interrupts in the early boot process.

remove_irq

LINUX

Kernel Hackers Manual July 2010

Name

`remove_irq` — free an interrupt

Synopsis

```
void remove_irq (unsigned int irq, struct irqaction * act);
```

Arguments

irq

Interrupt line to free

act

irqaction for the interrupt

Description

Used to remove interrupts statically setup by the early boot process.

free_irq

LINUX

Kernel Hackers Manual July 2010

Name

`free_irq` — free an interrupt allocated with `request_irq`

Synopsis

```
void free_irq (unsigned int irq, void * dev_id);
```

Arguments

irq

Interrupt line to free

dev_id

Device identity to free

Description

Remove an interrupt handler. The handler is removed and if the interrupt line is no longer in use by any driver it is disabled. On a shared IRQ the caller must ensure the interrupt is disabled on the card it drives before calling this function. The function does not return until any executing interrupts for this IRQ have completed.

This function must not be called from interrupt context.

request_threaded_irq

LINUX

Kernel Hackers Manual July 2010

Name

`request_threaded_irq` — allocate an interrupt line

Synopsis

```
int request_threaded_irq (unsigned int irq, irq_handler_t  
handler, irq_handler_t thread_fn, unsigned long irqflags,  
const char * devname, void * dev_id);
```

Arguments

irq

Interrupt line to allocate

handler

Function to be called when the IRQ occurs. Primary handler for threaded interrupts If NULL and `thread_fn != NULL` the default primary handler is installed

thread_fn

Function called from the irq handler thread If NULL, no irq thread is created

irqflags

Interrupt type flags

devname

An ascii name for the claiming device

dev_id

A cookie passed back to the handler function

Description

This call allocates interrupt resources and enables the interrupt line and IRQ handling. From the point this call is made your handler function may be invoked. Since your handler function must clear any interrupt the board raises, you must take care both to initialise your hardware and to set up the interrupt handler in the right order.

If you want to set up a threaded irq handler for your device then you need to supply *handler* and *thread_fn.handler* is still called in hard interrupt context and has to check whether the interrupt originates from the device. If yes it needs to disable the interrupt on the device and return `IRQ_WAKE_THREAD` which will wake up the handler thread and run *thread_fn*. This split handler design is necessary to support shared interrupts.

Dev_id must be globally unique. Normally the address of the device data structure is used as the cookie. Since the handler receives this value it makes sense to use it.

If your interrupt is shared you must pass a non NULL *dev_id* as this is required when freeing the interrupt.

Flags

`IRQF_SHARED` Interrupt is shared `IRQF_DISABLED` Disable local interrupts while processing `IRQF_SAMPLE_RANDOM` The interrupt can be used for entropy `IRQF_TRIGGER_*` Specify active edge(s) or level

set_irq_chip

LINUX

Kernel Hackers Manual July 2010

Name

`set_irq_chip` — set the irq chip for an irq

Synopsis

```
int set_irq_chip (unsigned int irq, struct irq_chip * chip);
```

Arguments

irq

irq number

chip

pointer to irq chip description structure

set_irq_type

LINUX

Kernel Hackers Manual July 2010

Name

`set_irq_type` — set the irq trigger type for an irq

Synopsis

```
int set_irq_type (unsigned int irq, unsigned int type);
```

Arguments

irq

irq number

type

IRQ_TYPE_{LEVEL,EDGE}_* value - see include/linux/irq.h

set_irq_data

LINUX

Kernel Hackers Manual July 2010

Name

`set_irq_data` — set irq type data for an irq

Synopsis

```
int set_irq_data (unsigned int irq, void * data);
```

Arguments

irq

Interrupt number

data

Pointer to interrupt specific data

Description

Set the hardware irq controller data for an irq

set_irq_chip_data

LINUX

Kernel Hackers Manual July 2010

Name

`set_irq_chip_data` — set irq chip data for an irq

Synopsis

```
int set_irq_chip_data (unsigned int irq, void * data);
```

Arguments

irq

Interrupt number

data

Pointer to chip specific data

Description

Set the hardware irq chip data for an irq

set_irq_nested_thread

LINUX

Kernel Hackers Manual July 2010

Name

`set_irq_nested_thread` — Set/Reset the IRQ_NESTED_THREAD flag of an irq

Synopsis

```
void set_irq_nested_thread (unsigned int irq, int nest);
```

Arguments

irq

Interrupt number

nest

0 to clear / 1 to set the IRQ_NESTED_THREAD flag

Description

The `IRQ_NESTED_THREAD` flag indicates that on `request_threaded_irq` no separate interrupt thread should be created for the irq as the handler are called nested in the context of a demultiplexing interrupt handler thread.

handle_level_irq

LINUX

Kernel Hackers Manual July 2010

Name

`handle_level_irq` — Level type irq handler

Synopsis

```
void handle_level_irq (unsigned int irq, struct irq_desc *  
desc);
```

Arguments

irq

the interrupt number

desc

the interrupt description structure for this irq

Description

Level type interrupts are active as long as the hardware line has the active level. This may require to mask the interrupt and unmask it after the associated handler has acknowledged the device, so the interrupt line is back to inactive.

Chapter 8. Public Functions Provided

Chapter 9. Internal Functions Provided

This chapter contains the autogenerated documentation of the internal functions.

handle_bad_irq

LINUX

Kernel Hackers Manual July 2010

Name

`handle_bad_irq` — handle spurious and unhandled irqs

Synopsis

```
void handle_bad_irq (unsigned int irq, struct irq_desc *  
desc);
```

Arguments

irq

the interrupt number

desc

description of the interrupt

Description

Handles spurious and unhandled IRQ's. It also prints a debugmessage.

handle_IRQ_event

LINUX

Kernel Hackers Manual July 2010

Name

handle_IRQ_event — irq action chain handler

Synopsis

```
irqreturn_t handle_IRQ_event (unsigned int irq, struct  
irqaction * action);
```

Arguments

irq

the interrupt number

action

the interrupt action chain for this irq

Description

Handles the action chain of an irq event

__do_IRQ

LINUX

Name

`__do_IRQ` — original all in one highlevel IRQ handler

Synopsis

```
unsigned int __do_IRQ (unsigned int irq);
```

Arguments

irq

the interrupt number

Description

`__do_IRQ` handles all normal device IRQ's (the special SMP cross-CPU interrupts have their own specific handlers).

This is the original x86 implementation which is used for every interrupt type.

dynamic_irq_init

LINUX

Name

`dynamic_irq_init` — initialize a dynamically allocated irq

Synopsis

```
void dynamic_irq_init (unsigned int irq);
```

Arguments

irq

irq number to initialize

dynamic_irq_init_keep_chip_data

LINUX

Kernel Hackers Manual July 2010

Name

`dynamic_irq_init_keep_chip_data` — initialize a dynamically allocated irq

Synopsis

```
void dynamic_irq_init_keep_chip_data (unsigned int irq);
```

Arguments

irq

irq number to initialize

Description

does not set `irq_to_desc(irq)->chip_data` to NULL

dynamic_irq_cleanup

LINUX

Kernel Hackers Manual July 2010

Name

`dynamic_irq_cleanup` — cleanup a dynamically allocated irq

Synopsis

```
void dynamic_irq_cleanup (unsigned int irq);
```

Arguments

irq

irq number to initialize

dynamic_irq_cleanup_keep_chip_data

LINUX

Name

`dynamic_irq_cleanup_keep_chip_data` — cleanup a dynamically allocated irq

Synopsis

```
void dynamic_irq_cleanup_keep_chip_data (unsigned int irq);
```

Arguments

irq

irq number to initialize

Description

does not set `irq_to_desc(irq)->chip_data` to NULL

set_irq_msi

LINUX

Name

`set_irq_msi` — set MSI descriptor data for an irq

Synopsis

```
int set_irq_msi (unsigned int irq, struct msi_desc * entry);
```

Arguments

irq

Interrupt number

entry

Pointer to MSI descriptor data

Description

Set the MSI descriptor entry for an irq

handle_simple_irq

LINUX

Kernel Hackers Manual July 2010

Name

`handle_simple_irq` — Simple and software-decoded IRQs.

Synopsis

```
void handle_simple_irq (unsigned int irq, struct irq_desc *  
desc);
```

Arguments

irq

the interrupt number

desc

the interrupt description structure for this irq

Description

Simple interrupts are either sent from a demultiplexing interrupt handler or come from hardware, where no interrupt hardware control is necessary.

Note

The caller is expected to handle the ack, clear, mask and unmask issues if necessary.

handle_fasteoi_irq

LINUX

Kernel Hackers Manual July 2010

Name

`handle_fasteoi_irq` — irq handler for transparent controllers

Synopsis

```
void handle_fasteoi_irq (unsigned int irq, struct irq_desc *  
desc);
```

Arguments

irq

the interrupt number

desc

the interrupt description structure for this irq

Only a single callback will be issued to the chip

an `->eoi` call when the interrupt has been serviced. This enables support for modern forms of interrupt handlers, which handle the flow details in hardware, transparently.

handle_edge_irq

LINUX

Kernel Hackers Manual July 2010

Name

`handle_edge_irq` — edge type IRQ handler

Synopsis

```
void handle_edge_irq (unsigned int irq, struct irq_desc *  
desc);
```

Arguments

irq

the interrupt number

desc

the interrupt description structure for this irq

Description

Interrupt occurs on the falling and/or rising edge of a hardware signal. The occurrence is latched into the irq controller hardware and must be acked in order to be reenabled. After the ack another interrupt can happen on the same source even before the first one is handled by the associated event handler. If this happens it might be necessary to disable (mask) the interrupt depending on the controller hardware. This requires to reenale the interrupt inside of the loop which handles the interrupts which have arrived while the handler was running. If all pending interrupts are handled, the loop is left.

handle_percpu_irq

LINUX

Kernel Hackers Manual July 2010

Name

handle_percpu_irq — Per CPU local irq handler

Synopsis

```
void handle_percpu_irq (unsigned int irq, struct irq_desc *  
desc);
```

Arguments

irq

the interrupt number

desc

the interrupt description structure for this irq

Description

Per CPU interrupts on SMP machines without locking requirements

Chapter 10. Credits

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